1. Downward self-reducibility. For a language $A$, define

$$A_{<n} = \{ x \in A : |x| < n \}.$$  

Language $A$ is said to be downward self-reducible if it is possible to determine in polynomial time if $x \in A$ using the results of queries of the form “$y \in A_{<|x|}$?” The queries may be adaptive, meaning that the polynomial time procedure may choose later queries depending on the results of earlier ones. Show that every downward self-reducible language is in $PSPACE$.

2. Show that one of the following inequalities must hold: $L \neq P$ or $P \neq PSPACE$. Note that both are believed to be true, and no one knows how to prove either one is true.

3. Show that logspace reductions are closed under composition. Then use the same ideas to prove that if language $A$ is $P$-complete, then $A \in L$ implies $L = P$.

4. Use a padding argument to show that if $L = P$ then $PSPACE = EXP$.

5. Prove that $SPACE(O(n)) \neq P$. (Note that while this is an interesting result, it doesn’t seem to shed any light on the major open questions $L \nsubseteq P$ and $P \nsubseteq PSPACE$). Hint: consider a language $A$ and a padded version of $A$. How are the two languages related with respect to space? How are they related with respect to time?