

4.2 Lecture 21 (19/Nov): $G(n, p)$ thresholds

4.2.1 Threshold for H as a subgraph in $G(n, p)$

Working with low moments of random variables can be incredibly effective, even when we are not specifically looking for limited-independence sample spaces. Here is a prototypical example. “When” does a specific, constant-sized graph H , show up as a subgraph of a random graph selected from the distribution $G(n, p)$? We have in mind that we are “turning the knob” on p . If H has any edges then when $p = 0$, with probability 1 there is no subgraph isomorphic to H . When $p = 1$, with probability 1 such subgraphs are everywhere². In between, for any finite n , the probability is some increasing function of p . But we won’t take n finite, we will take it tending to ∞ .

So the question is,³ can we identify a function $\pi(n)$ such that in the model $G(n, p(n))$, with $\llbracket H \rrbracket$ denoting the event that there is an H in the random graph G ,

(a) If $p(n) \in o(\pi(n))$, then $\lim_n \Pr(\llbracket H \rrbracket) = 0$.

(b) If $p(n) \in \omega(\pi(n))$, then $\lim_n \Pr(\llbracket H \rrbracket) = 1$.

Such a function $\pi(n)$ is known as the *threshold* for appearance of H . It follows from work of Bollobas and Thomason [19] that monotone events—events that must hold in G' if they hold in some $G \subseteq G'$ —always have a threshold function.

(A related but incomparable statement: for a monotone graph property, i.e., a monotone property invariant under vertex permutations, for any $\varepsilon > 0$ there is a $p(n)$ such that $\Pr_{p(n)}(\text{property}) \leq \varepsilon$ and $\Pr_{p(n)+O(1/\log n)}(\text{property}) \geq 1 - \varepsilon$. See [42].)

4.2.2 Most pairs independent: threshold for K_4 in $G(n, p)$

Let $S \subset \{1, \dots, n\}$, $|S| = 4$. Let X_S be the event that K_4 occurs as a subgraph of G at S —that is, when you look at those four vertices, all the edges between them are present. Conflating X_S with its indicator function and letting X be the number of K_4 ’s in G , we have

$$X = \sum_S X_S$$

and

$$E(X) = \binom{n}{4} p^6.$$

We are interested in $\Pr(X > 0)$. Let $\pi(n) = n^{-2/3}$.

(a) For $p(n) \in o(\pi(n))$, $E(X) \in o(1)$, so $\Pr(\llbracket K_4 \rrbracket) \in o(1)$ and therefore $\lim_n \Pr(\llbracket K_4 \rrbracket) = 0$.

(b) For $1 > p(n) \in \omega(\pi(n))$, $E(X) \in \omega(1)$. We’d like to conclude that likely $X > 0$ but we do not have enough information to justify this, as it could be that X is usually 0 and occasionally very large.⁴ We will exclude that possibility for K_4 by studying the next moment of the distribution.

Before carrying out this calculation, though, we have to make one important note. Since the event $\llbracket K_4 \rrbracket$ is monotone, $[p \leq p'] \Rightarrow [\Pr_{G(n,p)} \llbracket K_4 \rrbracket \leq \Pr_{G(n,p')} \llbracket K_4 \rrbracket]$. (An easy way to see this is by choosing reals iid uniformly in $[0, 1]$ at each edge, and placing the edge in the graph if the rv is above the p

²Today we focus on $H = K_4$, the 4-clique, but more generally this method will establish the probability of any fixed graph H occurring as a subgraph in G , that is, there is an injection of vertices of H into the vertices of G such that every edge of H is present in G . This is different from asking that H occur as an *induced* subgraph of G , which requires also that non-edges of H be non-edges in G . That is an interesting question too but different in an essential way: the event is not monotone in G .

³Recall $p(n) \in o(\pi(n))$ means that $\limsup p(n)/\pi(n) = 0$, and $p(n) \in \omega(\pi(n))$ means that $\limsup \pi(n)/p(n) = 0$.

⁴When we study not K_4 -subgraphs, but other subgraphs, this can really happen. We’ll discuss this below.

or p' threshold.) This means that it is enough to show that K_4 “shows up” slightly above π . This is useful because some of our calculations break down far above π , not because there is anything wrong with the underlying statement but because the inequalities we use are not strong enough to be useful there and a direct calculation would need to take account of further moments.

To simplify our remaining calculations, then, let

$$p = n^{-2/3}g(n), \quad \text{so } n^4p^6 = g^6$$

for any sufficiently small $g(n) \in \omega(1)$; we’ll see how this is helpful in the calculations.

By an earlier exercise,

$$\text{Var}(X) = \sum_S \text{Var}(X_S) + \sum_{S \neq T} \text{Cov}(X_S, X_T)$$

X_S is a coin (or Bernoulli rv) with $\Pr(X_S = 1) = p^6$. The variance of such an rv is $p^6(1 - p^6)$.

The covariance terms are more interesting.

1. If $|S \cap T| \leq 1$, no edges are shared, so the events are independent and $\text{Cov}(X_S, X_T) = 0$.
2. If $|S \cap T| = 2$, one edge is shared, and a total of 11 specific edges must be present for both cliques to be present. A simple way to bound the covariance is (since $E(X_S), E(X_T) \geq 0$) that $\text{Cov}(X_S, X_T) = E(X_S X_T) - E(X_S)E(X_T) \leq E(X_S X_T) = p^{11}$.
3. If $|S \cap T| = 3$, three edges are shared, and a specific 9 edges must be present for both cliques to be present. Similarly to the previous case, $\text{Cov}(X_S, X_T) \leq p^9$.

$$\begin{aligned} \text{Var}(X) &\leq \binom{n}{4} p^6 (1 - p^6) + \binom{n}{2,2,2} p^{11} + \binom{n}{3,1,1} p^9 \\ &\in O(n^4 p^6 + n^6 p^{11} + n^5 p^9) \\ &= O(g^6 n^{4-4} + g^{11} n^{6-22/3} + g^9 n^{5-6}) \quad \text{from } p = n^{-2/3} g(n) \\ &= O(g^6 + g^{11} n^{-4/3} + g^9 n^{-1}) \\ &= O(g^6) \quad \text{provided } g^5 \in O(n^{4/3}) \text{ and } g^3 \in O(n) \end{aligned} \tag{4.4}$$

This gives us the key piece of information. For $g \in \omega(1)$ but not *too* large, we have

$$\frac{\text{Var}(X)}{(E(X))^2} \in \frac{O(g^6)}{\Theta((n^4 p^6)^2)} = \frac{O(g^6)}{\Theta(g^{12})} = O(g^{-6}) \subseteq o(1)$$

and we have only to apply the Chebyshev inequality (Cor. 14) or better yet Paley-Zygmund 74 to conclude that $\Pr(X = 0) \in o(1)$ and so

$$\lim_n \Pr(\llbracket K_4 \rrbracket) = 1. \tag{4.5}$$

Since $\llbracket K_4 \rrbracket$ is a monotone event, (4.5) holds even for g above the range we needed for the calculation to hold. (Note, though, since there is so much “room” in the calculation, we could even have used the upper bound $O(g^{11})$ on 4.4, and not resorted to this monotonicity argument.)

Exercise: Show that the threshold for appearance, as a subgraph, of the graph with 5 edges and 4 vertices is $n^{-4/5}$.

Comment: For a general H the threshold for appearance of H in $G(n, p)$ as a subgraph is determined not by the ratio ρ_H of edges to vertices, but by the maximum of this ratio over induced subgraphs of H , call it $\rho_{\max H}$. We’ll see this on a problem set (and see [8]). If these numbers are different then above n^{-1/ρ_H} the expected number of H ’s starts tending to ∞ but almost certainly we have none; once we cross the higher threshold $n^{-1/\rho_{\max H}}$, there is an “explosion” of many of these subgraphs appearing. (They show up highly intersecting in the fewer copies of the critical induced subgraph.)