

Tracking Dynamic Networks under Sampling Constraints: Supporting Materials

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I. INTRODUCTION

This report contains supporting materials such as proofs, discussions, and additional numerical results, for [1]. See the original paper for terms and definitions.

II. PROOFS OF SELECTED THEOREMS

A. Relationship between Whittle's Index and Myopic Index

Proposition 2.1: For the link sampling problem, the Whittle's index (if exists) is no smaller than the myopic index, i.e., $W(x) \geq Y(x) \forall x \in [0, 1]$. Moreover, $W(x) \rightarrow Y(x)$ as $|p_{11} - p_{01}| \rightarrow 0$.

Proof: Due to the convexity of $V_{\beta,m}(x)$ [2], $V_{\beta,m}(\mathcal{T}(x)) \leq xV_{\beta,m}(p_{11}) + (1-x)V_{\beta,m}(p_{01})$. Since Whittle's index must satisfy $V_{\beta,W(x)}(x; U=0) = V_{\beta,W(x)}(x; U=1)$, plugging in the Bellman equations for $V_{\beta,m}(x; U=0)$ and $V_{\beta,m}(x; U=1)$ gives $W(x) + \max(x, 1-x) \geq 1$ and hence $W(x) \geq Y(x)$. As $p_{11} \rightarrow p_{01}$, equality will be achieved. ■

B. Threshold Structure of the Optimal Policy for Single-Armed Bandit with Subsidy

Lemma 2.2: The optimal policy for the single-armed bandit with subsidy m is a threshold policy: $\pi_m^*(x) = 1$ if and only if $\tau^-(m) < x < \tau^+(m)$, i.e., $\mathcal{P}(m) = [0, \tau^-(m)] \cup [\tau^+(m), 1]$.

Proof: Note that $V_{\beta,m}(x; U=1)$ is linear in x . By the convexity of the value function [2], we have that $V_{\beta,m}(x; U=0)$ is also convex in x . At $x=0$ or 1 , we have

$$\begin{cases} V_{\beta,m}(0; U=0) = m + 1 + \beta V_{\beta,m}(p_{01}), \\ V_{\beta,m}(1; U=0) = m + 1 + \beta V_{\beta,m}(p_{11}); \\ V_{\beta,m}(0; U=1) = 1 + \beta V_{\beta,m}(p_{01}), \\ V_{\beta,m}(1; U=1) = 1 + \beta V_{\beta,m}(p_{11}), \end{cases}$$

which implies that the endpoints of $V_{\beta,m}(x; U=0)$ are above, equal to, or below those of $V_{\beta,m}(x; U=1)$ for $m > 0$, $m = 0$, or $m < 0$, respectively, as illustrated in Fig. 1. Due to the convexity of $V_{\beta,m}(x; U=0)$, it must have at most two intersections with $V_{\beta,m}(x; U=1)$, denoted by $\tau^-(m)$ and $\tau^+(m)$; for cases without intersection, define $\tau^-(m) = \tau^+(m) \triangleq \tau^*$, where τ^* is the tangent point under a certain m_{\max} . Then as in Fig. 1, $V_{\beta,m}(x; U=1) > V_{\beta,m}(x; U=0)$ if and only if $x \in (\tau^-(m), \tau^+(m))$. ■

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Research was sponsored in part by the U.S. National Institute of Standards and Technology under Agreement Number 60NANB10D003 and a MURI funded through ARO Grant W911NF-06-1-0076.

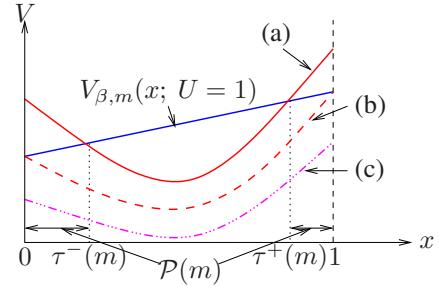


Fig. 1. Threshold structure of the optimal policy: Value function $V_{\beta,m}(x; U=0)$ for (a) $m > 0$, (b) $m = 0$, and (c) $m < 0$.

C. Pseudo-Linear Form of Value Function $V_{\beta,m}(x)$

Lemma 2.3: Given thresholds $\tau^-(m)$ and $\tau^+(m)$, define coefficients a_i, b_i ($i = 1, 2$) as in (3–6) and define functions

$$a(x) \triangleq \frac{1 - \beta^{L(x)}}{1 - \beta} + \beta^{L(x)+1} \mathcal{T}^{L(x)}(x) a_2 + \beta^{L(x)+1} (1 - \mathcal{T}^{L(x)}(x)) a_1, \quad (1)$$

$$b(x) \triangleq f(x; L(x)) + \beta^{L(x)} + \beta^{L(x)+1} \mathcal{T}^{L(x)}(x) b_2 + \beta^{L(x)+1} (1 - \mathcal{T}^{L(x)}(x)) b_1, \quad (2)$$

where $L(x) \triangleq \mathcal{L}(x; \tau^-(m), \tau^+(m))$. Then the value function is equal to $V_{\beta,m}(x) = a(x)m + b(x)$, with end values $V_{\beta,m}(p_{01}) = a_1m + b_1$, and $V_{\beta,m}(p_{11}) = a_2m + b_2$.

Proof: The linear forms of $V_{\beta,m}(p_{01})$ and $V_{\beta,m}(p_{11})$ are obtained by simply rewriting their expressions in (29, 30). Substituting them into (28) gives the linear form of $V_{\beta,m}(x)$. ■

D. Piecewise-Linear Property of $a(x), b(x)$

Proposition 2.4: The functions $a(x), b(x)$ defined in Lemma 3 of [1] are both piecewise-linear functions of x .

Proof: The proof is based on the piecewise-constant property of $L(x)$ (see Fig. 2). For fixed $L(x) \equiv l$, it is easy to see that $\mathcal{T}^{L(x)}(x)$ and $f(x; L(x))$ are both linear in x , which implies the linearity of $a(x)$ and $b(x)$. Thus, each constant piece of $L(x)$ corresponds to a linear piece of $a(x)$ and $b(x)$, respectively. ■

E. Monotonicity of $\tau^-(m), \tau^+(m)$

Lemma 2.5: The thresholds $\tau^-(m), \tau^+(m)$ are monotone increasing and decreasing, respectively, with m for $\beta \leq 0.5$.

Proof: It suffices to show ([3]) that for any given thresholds $(\tau^-(m'), \tau^+(m'))$ corresponding to some $m' \in [0, m_{\max}]$,

$$\frac{\partial}{\partial m} V_{\beta,m}(x; U=0) \geq \frac{\partial}{\partial m} V_{\beta,m}(x; U=1) \quad (7)$$

$$a_1 \triangleq \frac{(1 - \beta^{L_1+1} \mathcal{T}^{L_1}(p_{11})) \left(\frac{1-\beta^{L_2}}{1-\beta} \right) + \beta^{L_2+1} \mathcal{T}^{L_2}(p_{01}) \left(\frac{1-\beta^{L_1}}{1-\beta} \right)}{\eta}, \quad (3)$$

$$b_1 \triangleq \frac{(1 - \beta^{L_1+1} \mathcal{T}^{L_1}(p_{11}))(f(p_{01}; L_2) + \beta^{L_2}) + \beta^{L_2+1} \mathcal{T}^{L_2}(p_{01})(f(p_{11}; L_1) + \beta^{L_1})}{\eta}, \quad (4)$$

$$a_2 \triangleq \frac{(1 - \beta^{L_2+1}(1 - \mathcal{T}^{L_2}(p_{01}))) \left(\frac{1-\beta^{L_1}}{1-\beta} \right) + \beta^{L_1+1}(1 - \mathcal{T}^{L_1}(p_{11})) \left(\frac{1-\beta^{L_2}}{1-\beta} \right)}{\eta}, \quad (5)$$

$$b_2 \triangleq \frac{(1 - \beta^{L_2+1}(1 - \mathcal{T}^{L_2}(p_{01}))) (f(p_{11}; L_1) + \beta^{L_1}) + \beta^{L_1+1}(1 - \mathcal{T}^{L_1}(p_{11})) (f(p_{01}; L_2) + \beta^{L_2})}{\eta}. \quad (6)$$

for $x = \tau^-(m')$, $\tau^+(m')$, because this condition guarantees that for all $m > m'$, $V_{\beta,m}(x; U = 0) \geq V_{\beta,m}(x; U = 1)$ at $x = \tau^-(m')$ and $\tau^+(m')$, implying $\tau^-(m) \geq \tau^-(m')$, $\tau^+(m) \leq \tau^+(m')$. Next, by Lemma 2.3, we have

$$\begin{aligned} V_{\beta,m}(x; U = 0) &= m + \max(x, 1-x) + \beta a(\mathcal{T}(x))m \\ &\quad + \beta b(\mathcal{T}(x)), \\ V_{\beta,m}(x; U = 1) &= 1 + \beta x(a_2 m + b_2) + \beta(1-x)(a_1 m + b_1). \end{aligned}$$

Substituting these into (7) and noting that a_i, b_i and $a(\cdot), b(\cdot)$ are constants for fixed thresholds reduce (7) into

$$1 + \beta a(\mathcal{T}(x)) \geq \beta x a_2 + \beta(1-x) a_1, \quad x = \tau^-(m'), \tau^+(m'). \quad (8)$$

For $\beta \leq 1/2$, $1 + \beta a(\mathcal{T}(x)) \geq 1 \geq \beta/(1-\beta)$. Meanwhile, $a_1, a_2 \leq 1/(1-\beta)$ (since they are the discounted total passive time) implies $\beta/(1-\beta) \geq \beta x a_2 + \beta(1-x) a_1$, proving (8). ■

III. SUPPORTING STEPS IN COMPUTING WHITTLE'S INDEX

A. Computing Hitting Time $\mathcal{L}(x; c_1, c_2)$

For the ease of presentation, we introduce the following auxiliary functions:

$$g_1(y; x) \triangleq \frac{\log(\max(y - x_0, 0)) - \log|x - x_0|}{\log|p_{11} - p_{01}|}, \quad (9)$$

$$g_2(y; x) \triangleq \frac{\log(\max(x_0 - y, 0)) - \log|x - x_0|}{\log|p_{11} - p_{01}|}. \quad (10)$$

Then some calculation will show that for $p_{11} > p_{01}$,

$$\mathcal{L}(x; c_1, c_2) = \begin{cases} \min \mathbb{N} \cap (g_1(c_2; x), g_1(c_1; x)) & \text{if } x \geq x_0, \\ \min \mathbb{N} \cap (g_2(c_1; x), g_2(c_2; x)) & \text{if } x < x_0, \end{cases} \quad (11)$$

where \mathbb{N} denotes the set of nonnegative integers. For $p_{11} < p_{01}$, if $x \geq x_0$,

$$\mathcal{L}(x; c_1, c_2) = \min \left(\min \mathbb{N}_0 \cap (g_2(c_1; x), g_2(c_2; x)), \min \mathbb{N}_e \cap (g_1(c_2; x), g_1(c_1; x)) \right), \quad (12)$$

where \mathbb{N}_e is the set of nonnegative even numbers $(0, 2, \dots)$ and \mathbb{N}_o the set of nonnegative odd numbers $(1, 3, \dots)$. Similarly, if $x < x_0$,

$$\mathcal{L}(x; c_1, c_2) = \min \left(\min \mathbb{N}_o \cap (g_1(c_2; x), g_1(c_1; x)), \min \mathbb{N}_e \cap (g_2(c_1; x), g_2(c_2; x)) \right). \quad (13)$$

Sanity check: Consider the case $p_{11} > p_{01}$ (positively correlated arm). If $c_1 < x < c_2$, it is easy to see that

$\mathcal{L}(x; c_1, c_2) = 0$. Indeed, in (11), either $g_1(c_2; x) < 0$ and $g_1(c_1; x) > 0$ (if $x \geq x_0$), or $g_2(c_1; x) < 0$ and $g_2(c_2; x) > 0$ (if $x < x_0$), yielding $\mathcal{L}(x; c_1, c_2) = 0$. If $x, x_0 \geq c_2$ or $x, x_0 \leq c_1$, it is easy to see that $\mathcal{L}(x; c_1, c_2) = \infty$. Indeed, for $x, x_0 \geq c_2$, either $g_1(c_2; x) = \infty$ and $g_1(c_1; x) = \infty$ (if $x \geq x_0$), or $g_2(c_1; x) < 0$ and $g_2(c_2; x) < 0$ (if $x < x_0$); for $x, x_0 \leq c_1$, either $g_1(c_2; x) < 0$ and $g_1(c_1; x) < 0$ (if $x \geq x_0$), or $g_2(c_1; x) = \infty$ and $g_2(c_2; x) = \infty$ (if $x < x_0$), both yielding $\mathcal{L}(x; c_1, c_2) = \infty$ (define $(\infty, \infty) \triangleq \emptyset$ and $\min \emptyset \triangleq \infty$). Similar sanity check holds for $p_{11} < p_{01}$.

Due to the integral requirement, $\mathcal{L}(x; c_1, c_2)$ is always a piecewise-constant function of x for any c_1, c_2 and p_{01}, p_{11} , as illustrated in Fig. 2.

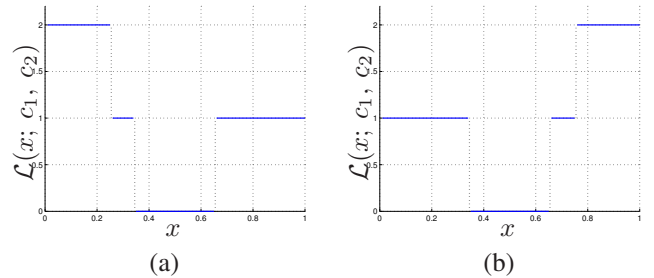


Fig. 2. Piece-wise constant property of $\mathcal{L}(x; c_1, c_2)$ ($c_1 = 0.35, c_2 = 0.65$): (a) $p_{01} = 0.25, p_{11} = 0.65$ (positively correlated); (b) $p_{01} = 0.65, p_{11} = 0.25$ (negatively correlated).

B. Computing Auxiliary Function $f(x; L)$

First of all, note that since L may be infinity, we cannot always compute $f(x; L)$ by the definition. Fortunately, due to its special structure, we can provide a closed-form solution as follows. The computation is based on the observation that $f(x; L)$ is a piecewise power series. We will treat positively-correlated and negatively-correlated arms separately.

For positively-correlated arms (*i.e.*, $p_{11} > p_{01}$), let $L_{1/2}$ denote the hitting time (*i.e.*, smallest l) for $\mathcal{T}^l(x)$ to cross $1/2$, if the crossing occurs within L steps. That is, $L_{1/2} \triangleq \min(L, \mathcal{L}(x; 1/2, 1))$ if $x \leq 1/2$, and $\min(L, \mathcal{L}(x; 0, 1/2))$ if $x > 1/2$. It is easy to see that

$$f(x; L) = \begin{cases} \sum_{i=0}^{L_{1/2}-1} \beta^i (1 - \mathcal{T}^i(x)) + \sum_{i=L_{1/2}}^{L-1} \beta^i \mathcal{T}^i(x) & \text{if } x \leq \frac{1}{2}, \\ \sum_{i=0}^{L_{1/2}-1} \beta^i \mathcal{T}^i(x) + \sum_{i=L_{1/2}}^{L-1} \beta^i (1 - \mathcal{T}^i(x)) & \text{o.w.} \end{cases}$$

By plugging in the expression of $\mathcal{T}^l(x)$, it can be shown that for $x \leq 1/2$,

$$f(x; L) = \frac{1 - \beta^{L_{1/2}}}{1 - \beta} - \frac{x_0(1 - 2\beta^{L_{1/2}} + \beta^L)}{1 - \beta} + \frac{(x_0 - x)(1 - 2(\beta(p_{11} - p_{01}))^{L_{1/2}} + (\beta(p_{11} - p_{01}))^L)}{1 - \beta(p_{11} - p_{01})}, \quad (14)$$

and for $x > 1/2$,

$$f(x; L) = \frac{\beta^{L_{1/2}} - \beta^L}{1 - \beta} + \frac{x_0(1 - 2\beta^{L_{1/2}} + \beta^L)}{1 - \beta} - \frac{(x_0 - x)(1 - 2(\beta(p_{11} - p_{01}))^{L_{1/2}} + (\beta(p_{11} - p_{01}))^L)}{1 - \beta(p_{11} - p_{01})}. \quad (15)$$

For negatively-correlated arms (*i.e.*, $p_{11} < p_{01}$), the even steps $\mathcal{T}^{2k}(x)$ and the odd steps $\mathcal{T}^{2k+1}(x)$ will converge toward x_0 from opposite directions. Let $\mathcal{K}(x; c_1, c_2) \triangleq \min\{k : \mathcal{T}^{2k} \in (c_1, c_2)\}$ denote the hitting time of (c_1, c_2) from x by taking two steps at a time, and $\mathcal{K}_{1/2}(x)$ the number of step pairs needed to first cross $1/2$ starting from x , *i.e.*, $\mathcal{K}_{1/2}(x) \triangleq \mathcal{K}(x; 1/2, 1)$ if $x \leq 1/2$, and $\mathcal{K}_{1/2}(x) \triangleq \mathcal{K}(x; 0, 1/2)$ if $x > 1/2$. Define $K_{1/2} \triangleq \min(\mathcal{K}_{1/2}(x), \lfloor (L-1)/2 \rfloor + 1)$, and $K'_{1/2} \triangleq \min(\mathcal{K}_{1/2}(\mathcal{T}(x)), \lfloor (L-2)/2 \rfloor + 1)$. Note that $\mathcal{K}(x; c_1, c_2)$ (thus $K_{1/2}, K'_{1/2}$) can be computed similarly as $\mathcal{L}(x; c_1, c_2)$ (see Section III-A). We can write $f(x; L)$ as

$$f(x; L) = \sum_{k=0}^{K_{1/2}-1} \beta^{2k} \max(\mathcal{T}^{2k}(x), 1 - \mathcal{T}^{2k}(x)) \quad (16)$$

$$+ \sum_{k=K_{1/2}}^{\lfloor \frac{L-1}{2} \rfloor} \beta^{2k} \max(\mathcal{T}^{2k}(x), 1 - \mathcal{T}^{2k}(x)) \quad (17)$$

$$+ \sum_{k=0}^{K'_{1/2}-1} \beta^{2k+1} \max(\mathcal{T}^{2k+1}(x), 1 - \mathcal{T}^{2k+1}(x)) \quad (18)$$

$$+ \sum_{k=K'_{1/2}}^{\lfloor \frac{L-2}{2} \rfloor} \beta^{2k+1} \max(\mathcal{T}^{2k+1}(x), 1 - \mathcal{T}^{2k+1}(x)). \quad (19)$$

This decomposition guarantees that (16) is on the same side of $1/2$ as x , (17) on the other side, (18) on the same side as $\mathcal{T}(x)$, and (19) on the other side.

We now calculate (16–19) by cases. If $x \leq 1/2$, then (16) is equal to $\sum_{k=0}^{K_{1/2}-1} \beta^{2k}(1 - \mathcal{T}^{2k}(x))$ and (17) to

$$\sum_{k=K_{1/2}}^{\lfloor \frac{L-1}{2} \rfloor} \beta^{2k} \mathcal{T}^{2k}(x). \text{ Calculation will yield the closed-form results as in (20–21). Otherwise (i.e., } x > 1/2), (16) \text{ becomes}$$

$$\sum_{k=0}^{K_{1/2}-1} \beta^{2k} \mathcal{T}^{2k}(x) \text{ and (17) becomes } \sum_{k=K_{1/2}}^{\lfloor \frac{L-1}{2} \rfloor} \beta^{2k}(1 - \mathcal{T}^{2k}(x)), \text{ which yield (22–23). Similarly, if } \mathcal{T}(x) \leq 1/2, \text{ then (18)}$$

$$\text{becomes } \sum_{k=0}^{K'_{1/2}-1} \beta^{2k+1}(1 - \mathcal{T}^{2k+1}(x)) \text{ and (19) becomes}$$

$$\sum_{k=K'_{1/2}}^{\lfloor \frac{L-2}{2} \rfloor} \beta^{2k+1} \mathcal{T}^{2k+1}(x), \text{ which gives the results in (24–25).}$$

Otherwise (*i.e.*, $\mathcal{T}(x) > 1/2$), (18) is $\sum_{k=0}^{K'_{1/2}-1} \beta^{2k+1} \mathcal{T}^{2k+1}(x)$

and (19) is $\sum_{k=K'_{1/2}}^{\lfloor \frac{L-2}{2} \rfloor} \beta^{2k+1}(1 - \mathcal{T}^{2k+1}(x))$, yielding (26–27).

C. Computing Value Function $V_{\beta,m}(x)$

It is shown in [1] that given m , $\tau^-(m)$, and $\tau^+(m)$, we can compute the value function of the single-armed bandit with subsidy m by

$$V_{\beta,m}(x) = \frac{(1 - \beta^L)m}{1 - \beta} + f(x; L) + \beta^L + \beta^{L+1} \mathcal{T}^L(x) V_{\beta,m}(p_{11}) + \beta^{L+1}(1 - \mathcal{T}^L(x)) V_{\beta,m}(p_{01}). \quad (28)$$

The only unknowns left are $V_{\beta,m}(p_{11}), V_{\beta,m}(p_{01})$. Note that $x = p_{11}$ or p_{01} should also satisfy (28), giving us two equations with two unknowns. Solving these equations yields the results in (29, 30).

IV. ADDITIONAL NUMERICAL RESULTS

We first verify the properties of $a(x), b(x)$ given in Proposition 2.4. As shown in Fig. 3, $a(x)$ and $b(x)$ are indeed piecewise-linear functions of x .

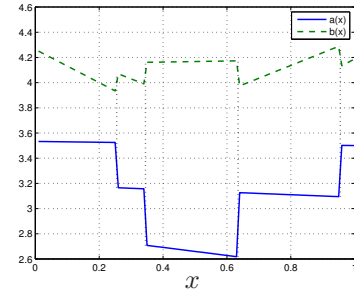


Fig. 3. Coefficients $a(x), b(x)$ vs. x ($\beta = 0.8, p_{01} = 0.25, p_{11} = 0.65, m = 0.4039, \tau^-(m) = 0.35, \tau^+(m) = 0.6329$).

We then verify the convexity of $V_{\beta,m}(x)$ with respect to m , which is needed to ensure that the performance upper bounds derived in [1] are well-defined and the associated subsidies are unique. We plot the value function $V_{\beta,m}(x_0)$ (with the steady state as the initial state) for the single-armed bandit as a function of subsidy m under positive and negative correlation, respectively, as shown in Fig. 4. In both cases, $V_{\beta,m}(x)$ is a monotone increasing, convex function of m . This observation holds even if we vary the parameters (not shown). Therefore, the expressions within the minimization of the bounds are convex in m (or \mathbf{m}), and hence the bounds are well-defined and the dual variables (subsidies) achieving them are unique.

Next, we compare the Whittle's policy and the myopic policy for the single-armed bandit with subsidy m , as shown in Fig. 5. We see that the two policies can behave differently even for a single arm; note that by the threshold structure in Lemma 2.2, the Whittle's index policy is also the optimal policy in the single-armed case. For large $|p_{11} - p_{01}|$, the myopic policy is strictly suboptimal (Fig. 5 (a)), and this holds over a range of subsidies (Fig. 5 (b)).

Finally, we verify the actual performance of the proposed policies measured by the total reward without discount:

$$x \leq \frac{1}{2} : (16) = \frac{(1-x_0)(1-\beta^{2K_{1/2}})}{1-\beta^2} + \frac{(x_0-x)[1-(\beta^2(p_{11}-p_{01})^2)^{K_{1/2}}]}{1-\beta^2(p_{11}-p_{01})^2}, \quad (20)$$

$$(17) = \frac{x_0(\beta^{2K_{1/2}} - \beta^{2(\lfloor \frac{L-1}{2} \rfloor + 1)})}{1-\beta^2} - \frac{(x_0-x)[(\beta^2(p_{11}-p_{01})^2)^{K_{1/2}} - (\beta^2(p_{11}-p_{01})^2)^{\lfloor \frac{L-1}{2} \rfloor + 1}]}{1-\beta^2(p_{11}-p_{01})^2}; \quad (21)$$

$$x > \frac{1}{2} : (16) = \frac{x_0(1-\beta^{2K_{1/2}})}{1-\beta^2} - \frac{(x_0-x)[1-(\beta^2(p_{11}-p_{01})^2)^{K_{1/2}}]}{1-\beta^2(p_{11}-p_{01})^2}, \quad (22)$$

$$(17) = \frac{(1-x_0)(\beta^{2K_{1/2}} - \beta^{2(\lfloor \frac{L-1}{2} \rfloor + 1)})}{1-\beta^2} + \frac{(x_0-x)[(\beta^2(p_{11}-p_{01})^2)^{K_{1/2}} - (\beta^2(p_{11}-p_{01})^2)^{\lfloor \frac{L-1}{2} \rfloor + 1}]}{1-\beta^2(p_{11}-p_{01})^2}. \quad (23)$$

$$\mathcal{T}(x) \leq \frac{1}{2} : (18) = \frac{(1-x_0)\beta(1-\beta^{2K'_{1/2}})}{1-\beta^2} + \frac{(x_0-x)\beta(p_{11}-p_{01})[1-(\beta^2(p_{11}-p_{01})^2)^{K'_{1/2}}]}{1-\beta^2(p_{11}-p_{01})^2}, \quad (24)$$

$$(19) = \frac{x_0\beta(\beta^{2K'_{1/2}} - \beta^{2(\lfloor \frac{L-2}{2} \rfloor + 1)})}{1-\beta^2} - \frac{(x_0-x)\beta(p_{11}-p_{01})[(\beta^2(p_{11}-p_{01})^2)^{K'_{1/2}} - (\beta^2(p_{11}-p_{01})^2)^{\lfloor \frac{L-2}{2} \rfloor + 1}]}{1-\beta^2(p_{11}-p_{01})^2}; \quad (25)$$

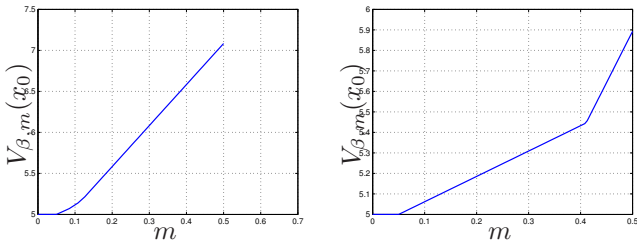
$$\mathcal{T}(x) > \frac{1}{2} : (18) = \frac{x_0\beta(1-\beta^{2K'_{1/2}})}{1-\beta^2} - \frac{(x_0-x)\beta(p_{11}-p_{01})[1-(\beta^2(p_{11}-p_{01})^2)^{K'_{1/2}}]}{1-\beta^2(p_{11}-p_{01})^2}, \quad (26)$$

$$(19) = \frac{(1-x_0)\beta(\beta^{2K'_{1/2}} - \beta^{2(\lfloor \frac{L-2}{2} \rfloor + 1)})}{1-\beta^2} + \frac{(x_0-x)\beta(p_{11}-p_{01})[(\beta^2(p_{11}-p_{01})^2)^{K'_{1/2}} - (\beta^2(p_{11}-p_{01})^2)^{\lfloor \frac{L-2}{2} \rfloor + 1}]}{1-\beta^2(p_{11}-p_{01})^2}. \quad (27)$$

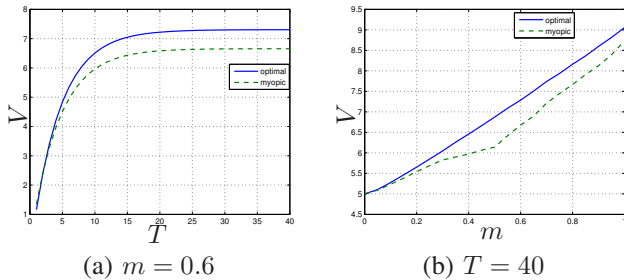
$$V_{\beta,m}(p_{01}) = \frac{(1-\beta^{L_1+1}\mathcal{T}^{L_1}(p_{11}))v_2 + \beta^{L_2+1}\mathcal{T}^{L_2}(p_{01})v_1}{\eta}, \quad (29)$$

$$V_{\beta,m}(p_{11}) = \frac{(1-\beta^{L_2+1}(1-\mathcal{T}^{L_2}(p_{01})))v_1 + \beta^{L_1+1}(1-\mathcal{T}^{L_1}(p_{11}))v_2}{\eta}, \quad (30)$$

where $L_1 \triangleq \mathcal{L}(p_{11}; \tau^-(m), \tau^+(m))$, $L_2 \triangleq \mathcal{L}(p_{01}; \tau^-(m), \tau^+(m))$, $v_1 \triangleq \frac{(1-\beta^{L_1})m}{1-\beta} + f(p_{11}; L_1) + \beta^{L_1}$, $v_2 \triangleq \frac{(1-\beta^{L_2})m}{1-\beta} + f(p_{01}; L_2) + \beta^{L_2}$, and $\eta \triangleq (1-\beta^{L_1+1}\mathcal{T}^{L_1}(p_{11}))(1-\beta^{L_2+1}(1-\mathcal{T}^{L_2}(p_{01}))) - \beta^{L_1+L_2+2}(1-\mathcal{T}^{L_1}(p_{11}))\mathcal{T}^{L_2}(p_{01})$.



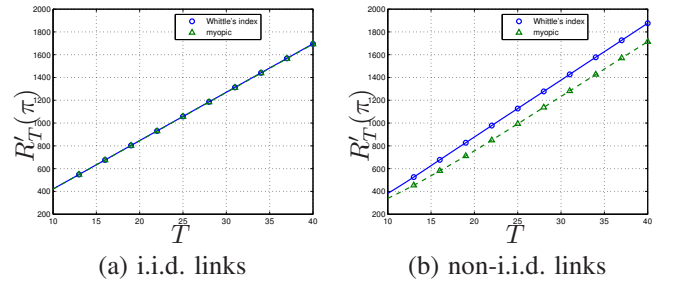
(a) positively-correlated arm (b) negatively-correlated arm
Fig. 4. $V_{\beta,m}(x_0)$ vs. m ($\beta = 0.8$): (a) $p_{01} = 0.05$, $p_{11} = 0.45$; (b) $p_{01} = 0.45$, $p_{11} = 0.05$.



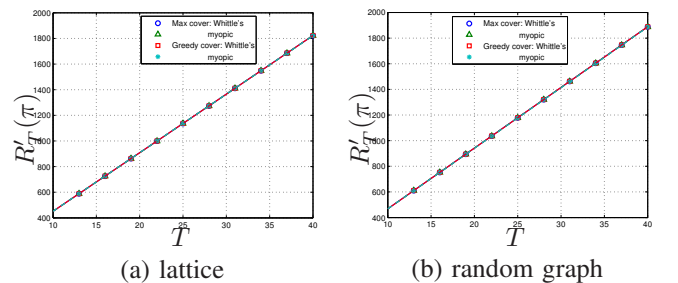
(a) $m = 0.6$ (b) $T = 40$
Fig. 5. Values of the optimal (Whittle's index) policy and the myopic policy for the single-armed bandit with subsidy ($\beta = 0.8$, $p_{01} = 0.01$, $p_{11} = 0.99$, 100 Monte Carlo runs).

$R'_T(\pi) := \mathbb{E}_\pi[\sum_{t=1}^T R(t; \pi)]$, i.e., the expected number of times until time step T that any link is tracked correctly. Under the same settings as in Section VI.C of [1], we plot the actual

performance in Fig. 6–8. We see that the comparison results between the myopic policy and the (extended) Whittle's index policy in [1] carry through to the actual performance as well.



(a) i.i.d. links (b) non-i.i.d. links
Fig. 6. Link sampling ($\beta = 0.8$, $M = 60$, $K = 3$, $T = 40$, $\mathbf{x}(1) = \mathbf{x}_0$, 100 Monte Carlo runs): (a) $p_{01} = 0.2$, $p_{11} = 0.9$; (b) $p_{01} = 0.999$ for fast links and 0.001 for slow links, $p_{11} = 1 - p_{01}$.



(a) lattice (b) random graph
Fig. 7. Node sampling: i.i.d. links ($\beta = 0.8$, $N = 36$, $M = 60$, $K = 2$, $T = 40$, $p_{01} = 0.2$, $p_{11} = 0.9$, $\mathbf{x}(1) = \mathbf{x}_0$, 100 Monte Carlo runs).

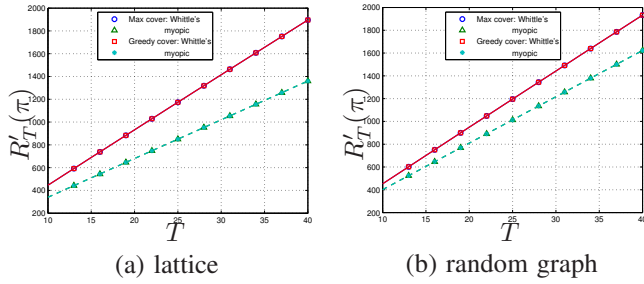


Fig. 8. Node sampling: non-i.i.d. links ($p_{01} = 0.5$ for fast links and 0.001 for slow links, $p_{11} = 1 - p_{01}$, rest as in Fig. 7).

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