Algorithmic Models for Social Network Phenomena

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Networks as Phenomena



- Complex networks as phenomena, not just designed artifacts.
- What recurring patterns emerge, why are they there, and what are the consequences for computing and information systems?

Social and Technological Networks



Social networks: friendships, contacts, collaboration, influence, organizational structure, economic institutions.

- Social and technological networks are intertwined: Web content, blogging, e-mail/IM, MySpace/Facebook/...
- New technologies change our patterns of social interaction.
- Collecting social data at unprecedented scale and resolution.

Rich Social Network Data

Traditional obstacle: Can only choose 2 of 3.

- Large-scale
- Realistic
- Completely mapped



Two lines of research, looking for a meeting point.

- Social scientists engaged in detailed study of small datasets, concerned with social outcomes.
- Computer scientists discovering properties of massive network datasets that were invisible at smaller scales.

Modeling Complex Networks

We want Kepler's Laws of Motion for the Web.

Mike Steuerwalt,
 NSF KDI Workshop, 1998



Opportunity for deeper understanding of information networks and social processes, informed by theoretical models and rich data.

- Mathematical / algorithmic models form the vocabulary for expressing complex social-science questions on complex network data.
- Payoffs from the introduction of an algorithmic perspective into the social sciences.

Overview

(1) Small-world networks and decentralized search

- Stylized models expose basic patterns.
- Identifying the patterns in large-scale data.
- (2) A problem that is less well understood at a large scale: diffusion and cascading behavior in social networks
 - The way in which new practices, ideas, and behaviors spread through social networks like epidemics.
 - Models from discrete probability, data from on-line communities, open questions in relating them.
- (3) Privacy and anonymity in on-line data.
 - The perils in using anonymized social network data.
 - Attacks on anonymized networks using small identifiable subgraphs.

Milgram's small-world experiment (1967)

Choose a target in Boston, starters in Nebraska.
A letter begins at each starter, must be passed between personal acquaintances until target is reached.
Six steps on average → six degrees of separation.

- Routing in a (social) network: When is local information sufficient? [Kleinberg 2000]
- Variation on network model of Watts and Strogatz [1998].
- Add edges to lattice: u links to v with probability $d(u, v)^{-\alpha}$.



Small-World Models

- Optimal exponent $\alpha = 2$: yields routing time $\sim c \log^2 n$.
- All other exponents yield $\sim n^{\varepsilon}$ for some $\varepsilon > 0$.



- Generalizations to random networks on different "scaffolds":
 - Trees, set systems, low-dimensional metrics [Kleinberg '01, Watts-Dodds-Newman '02, Slivkins '05,

Fraigniaud-Lebhar-Lotker '06, Abraham-Gavoille '06]

• Relation to long-range percolation, structured random graphs

- [Newman-Schulman'86, Aizenman-Chayes-Chayes-Newman'88, Bollobás-Chung '88, Benjamini-Berger '01, Coppersmith-Gamarnik-Sviridenko '02, Biskup '04, Berger '06]
- Connections to peer-to-peer algorithms
 - [Kempe-Kleinberg-Demers '01, Malkhi-Naor-Ratajczak '02, Aspnes-Diamadi-Shah '02, Zhang-Goel-Govindan '02, Manku-Bawa-Raghavan '03, Li et al. '05]

Social Network Data

- [Adamic-Adar 2003]: social network on 436 HP Labs researchers.
- Joined pairs who exchanged ≥ 6 e-mails (each way).



- Compared to "group-based" model [Kleinberg 2001]
 - Probability of link (v, w) prop. to g(v, w)^{-α}, where g(v, w) is size of smallest group containing v and w.
 - $\alpha = 1$ gives optimal search performance.
- In HP Labs, groups defined by sub-trees of hierarchy.
- Links scaled as $g^{-3/4}$.

Geographic Data: LiveJournal



Liben-Nowell, Kumar, Novak, Raghavan, Tomkins (2005) studied LiveJournal, an on-line blogging community with friendship links.

- Large-scale social network with geographical embedding:
 - 500,000 members with U.S. Zip codes, 4 million links.
- Analyzed how friendship probability decreases with distance.
- Difficulty: non-uniform population density makes simple lattice models hard to apply.

LiveJournal: Rank-Based Friendship



Rank-based friendship: <u>rank</u> of w with respect to v is number of people x such that d(v, x) < d(v, w).

- Decentralized search with (essentially) arbitrary population density, when link probability proportional to $rank^{-\beta}$.
- (LKNRT'05): Efficient routing when $\beta = 1$, i.e. 1/rank.
- Generalization of lattice result (diff. from set systems).

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Punchline: LiveJournal friendships approximate 1/rank.

Diffusion in Social Networks



So far: focused search in a social network.

Now switch to diffusion, another fundamental social processs: Behaviors that cascade from node to node like an epidemic.

- News, opinions, rumors, fads, urban legends, ...
- Word-of-mouth effects in marketing, rise of new products.
- Changes in social priorities: smoking, recycling, ...
- Saturation news coverage; topic diffusion among bloggers.
- Localized collective action: riots, walkouts

Diffusion Curves

Basis for models: Probability of adopting new behavior depends on number of friends who have adopted.

• Bass 1969; Granovetter 1978; Schelling 1978



Key issue: qualitative shape of the diffusion curves.

• Diminishing returns? Critical mass?

From individual-level model, can build network-level model:

• Run dynamics of contagion forward from initial "seed set."

Finding the Most Influential Set



An algorithmic question [Domingos-Richardson 2001]:

• If we can "seed" the new behavior at k nodes, and want to maximize the eventual spread, whom should we choose?

Computational complexity depends on diffusion curves.

- Highly inapproximable with critical mass.
- With diminishing returns: constant-factor approximation [Kempe-Kleinberg-Tardos 2003, 2005; Mossel-Roch 2007]

Diffusion Curves







Editing a Wikipedia article [Huttenlocher et al. '07] Triadic closure in e-mail [Kossinets-Watts '06]

Toward a Notion of "Life Cycles"

How does a group's tendency to grow depend on its structural properties? [Backstrom et al. 2006]





- Define clustering = # triangles / # open triads.
- Look at growth from t_1 to t_2 as function of clustering.
- Groups with large clustering grow slower.
- Yet individuals are more likely to join when their friends in a group all know each other.



Diffusion in Computing and Information

- Diffusion of Topics [Gruhl et al 2004, Adar et al 2004]
 - News stories cascade through networks of bloggers and media
 - How should we track stories and rank news sources?
 - A taxonomy of sources: discoverers, amplifiers, reshapers, ...
- Building diffusion into the design of social media [Leskovec-Adamic-Huberman 2006, Kleinberg-Raghavan 2005]
 - Incentives to propagate interesting recommendations along social network links.
 - Simple markets based on question-answering and information-seeking.
- Predictive frameworks for diffusion
 - Machine learning models for the growth of communities [Backstrom et al. 2006]
 - Is a new idea's rise to success inherently unpredictable? [Salganik-Dodds-Watts 2006]

The Perils of Anonymized Data

Can accomplish a lot with public social network data. But many interesting questions arise in private data:

- E.g. E-mail, IM, voice, members-only communities.
- Standard approach to protecting the data: anonymize, replacing name at each node by a random string.
- After doing this, is it safe to release?

With more detailed data, anonymization has run into trouble:

- Identifying on-line pseudonyms by textual analysis [Novak-Raghavan-Tomkins 2004]
- De-anonymizing Netflix ratings via time series [Narayanan-Shmatikov 2006]
- The AOL query logs ["This was a screw-up, and we're angry and upset about it." —AOL press release, 7 August 2006]

But what about just the unlabeled nodes and edges of a social network?



Scenario from Backstrom-Dwork-Kleinberg 2007:

Suppose a big company were going to release an anonymized communication graph on 100 million users.



An attacker chooses a small set of *b* user accounts to "target": Goal is to learn edge relations among them.



Before dataset is released:

- Create a small set of k fake new accounts, with links among them, forming a subgraph H.
- Attach this new subgraph H to targeted accounts.



When anonymized dataset is released, need to find H.

Why couldn't there be many copies of *H* in the dataset? Isn't subgraph isomorphism supposed to be a hard problem?



If H is random and of size $(2 + \varepsilon) \log n$, then:

It's unique with high probability

(cf. Erdös's (1947) non-constructive Ramsey bound).

Brute-force search tree for H has near-linear size, since H is small and random.



Once *H* is found:

Can easily find the targeted nodes by following edges from H.

First version of the attack:

- Create H on (2 + ε) log n nodes.
 Can compromise Θ(log² n) targeted nodes.
- In experiments on 4.4 million-node LiveJournal graph, 7-node graph H can compromise 70 targeted nodes (and hence ~ 2400 edge relations).

Second version of the attack:

- Create *H* on $c\sqrt{\log n}$ nodes. Can compromise $(\frac{1}{2} - \varepsilon)c\sqrt{\log n}$ targeted nodes.
- Reconstruct from Gomory-Hu tree: break apart *G* along small cuts; find *H* as a "contiguous" piece.

Passive attacks:

• In LiveJournal graph: with reasonable probability, you and 6 of your friends chosen at random can carry out the first attack, compromising about 10 users.

What's the conclusion from this?

- Doesn't apply to social network data that's already public; orthogonal to issues of legal/contractual safeguards.
- But widespread release of an anonymized social network? Danger: you don't what someone's hidden in there. (And passive attacks don't even require advance planning.)
- Interesting direction: privacy-preserving mechanisms for making social network data accessible.
 - May be difficult to obfuscate network effectively (e.g. [Dinur-Nissim 2003, Dwork-McSherry-Talwar 2007])
 - Interactive mechanisms for network data may be possible (e.g. [Dwork-McSherry-Nissim-Smith 2006])

Reflections: Toward a Model of You

Further direction: from populations to individuals

- Distributions over millions of people leave open several possibilities:
 - Each individual personally follows (a version of) the distribution, or
 - Individual are highly diverse, and the distribution only appears in aggregate
- Recent studies suggests that sometimes the first option may in fact be true.

Example: what is the probability that you answer a piece of e-mail *t* days after receipt (conditioned on answering at all)?

 Recent theories suggest t^{-1.5} with exponential cut-off [Barabasi 2005]



Reflections: Interacting in the On-Line World

MySpace is doubly awkward because it makes public what should be private. It doesn't just create social networks, it anatomizes them. It spreads them out like a digestive tract on the autopsy table. You can see what's connected to what, who's connected to whom.

- Toronto Globe and Mail, June 2006.
- Social networks implicit for millenia are increasingly being recorded at arbitrary resolution and browsable in our information systems.
- Your software has a trace of your activities resolved to the second — and increasingly knows more about your behavior than you do.
- Trade-offs between rich data and individual privacy will remain an issue.
- Models based on algorithmic ideas will be crucial in understanding these developments.